



# SPETABARU: A Task-based Runtime System with Speculative Execution Capability

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# SPETABARU: A Task-based Runtime System with Speculative Execution Capability

Bérenger Bramas

Inria Nancy - Grand Est / CAMUS Team



# Summary

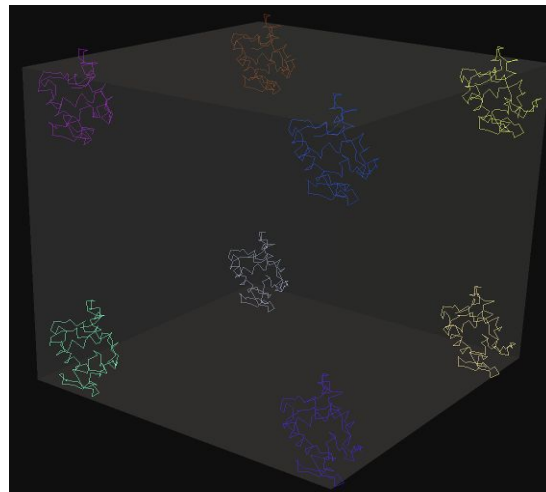
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- Background & motivation
  - Monte Carlo (MC), Replica Exchange Monte Carlo (REMC)
  - Task-based parallelization (STF)
- Speculative execution in task-based runtime systems
  - Introduction
  - Challenges
  - Some solutions
- SPETABARU: SPEculative TAsk BAsed RUnTime
  - Illustration
  - Performance study
- Conclusion & perspective

## Motivation - Biophysic application

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- Can be viewed as a molecular dynamic code
- A protein is composed of domains/polypeptides/chains of amino acids
  - Each domain is composed of beads/particles
  - A domain can “move/rotate/shift”
  - The global energy is computed with  $N^2$  particles-particles interactions
- Too many configurations -> MC



# Motivation - MC

Write(d1)

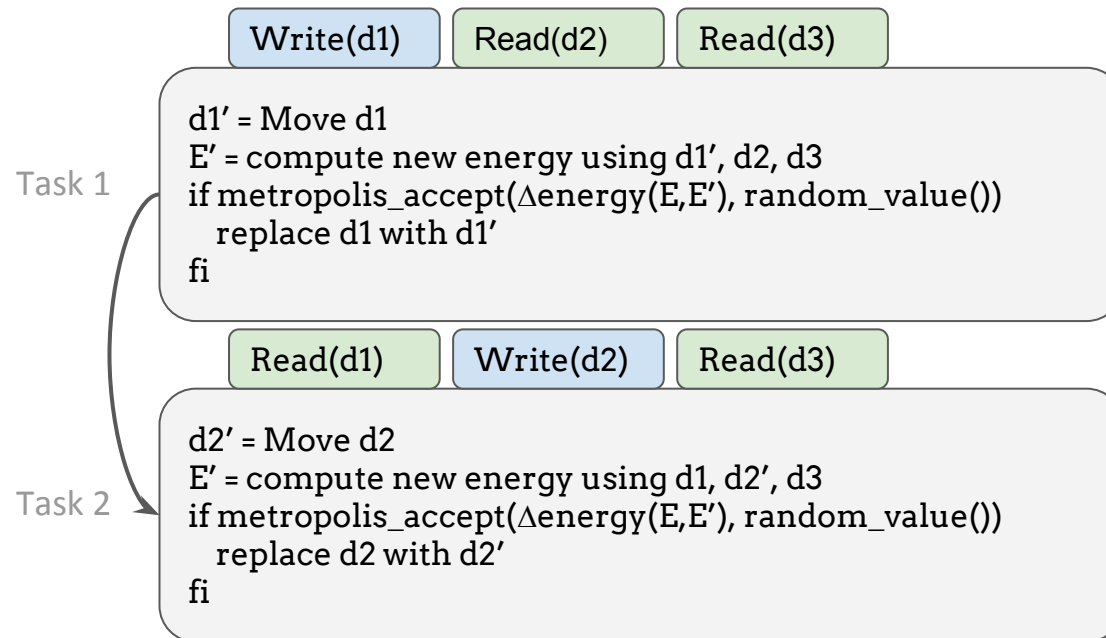
Read(d2)

Read(d3)

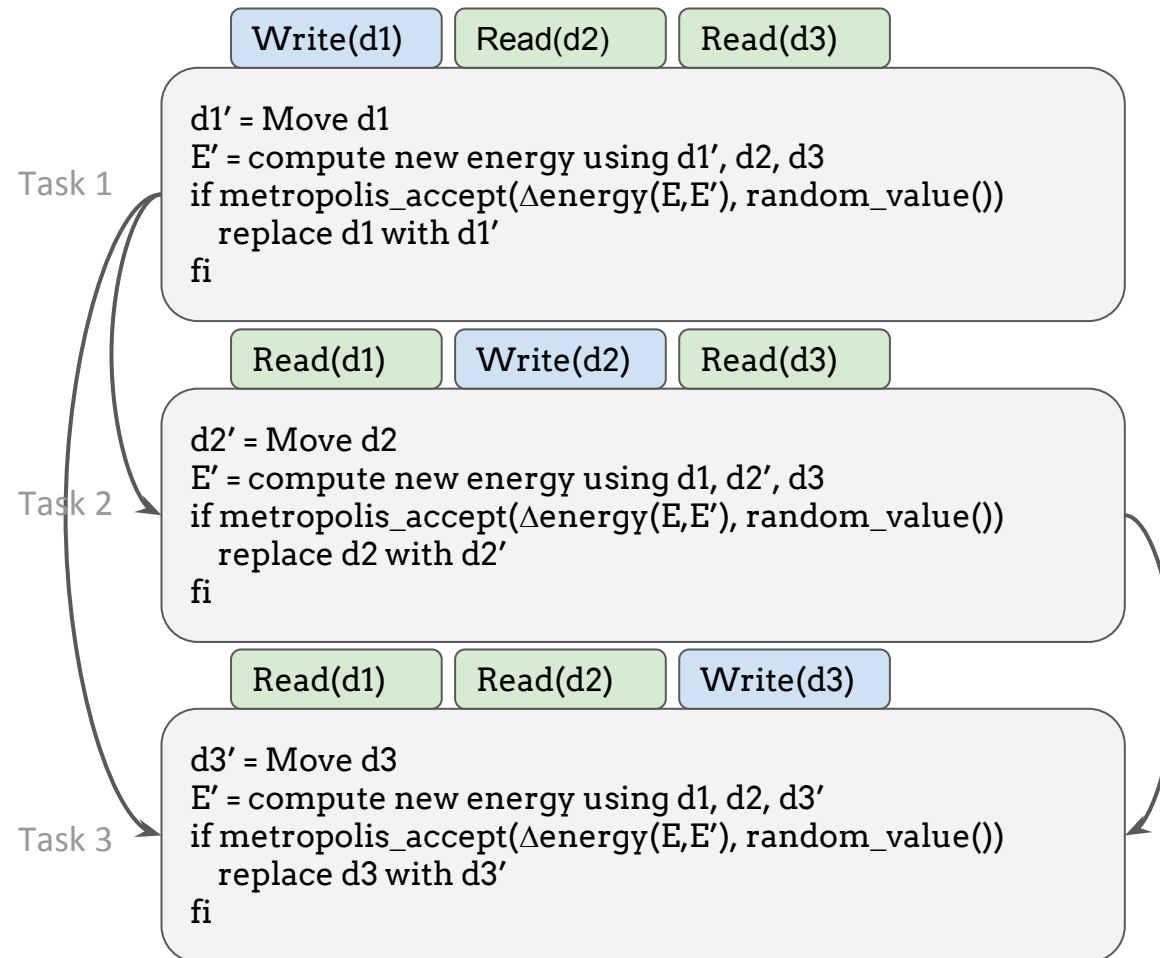
Task 1

```
d1' = Move d1
E' = compute new energy using d1', d2, d3
if metropolis_accept( $\Delta$ energy(E,E'), random_value())
  replace d1 with d1'
fi
```

## Motivation - MC

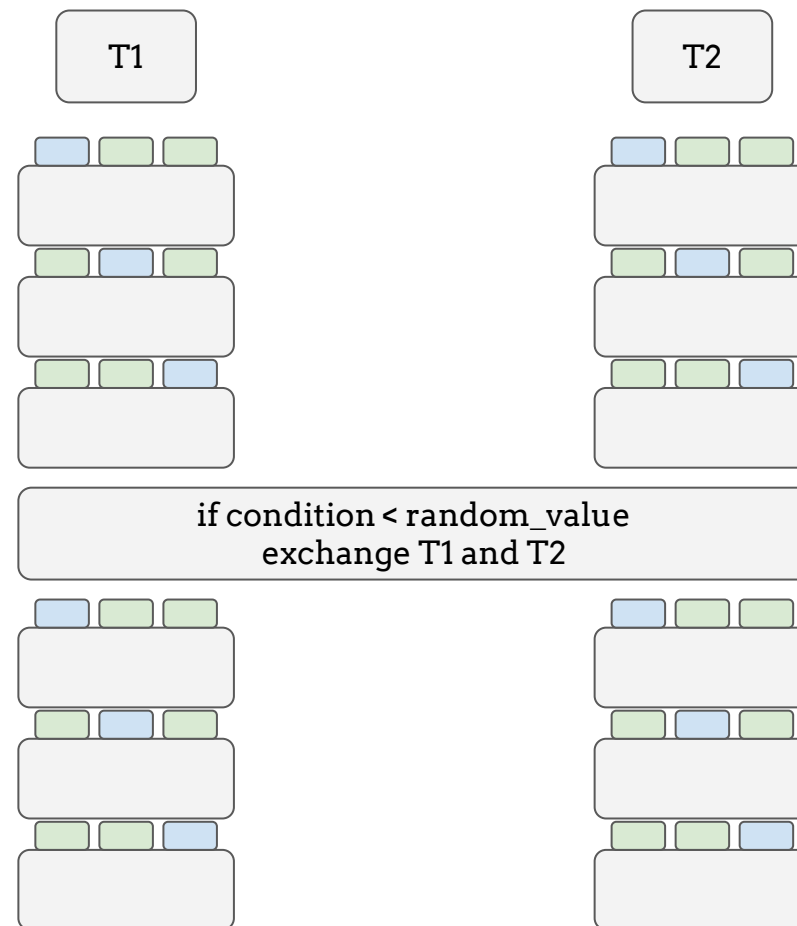


# Motivation - MC



## Motivation - REMC

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# Problem

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- Limited degree of parallelism
  - Each iteration depends on the previous ones
  - Possible exchange of temperature = dependency
- Some tasks access the data in “write” but it is unknown at task insertion time if they are really going to modify the data or not

Write(d1)

Read(d2)

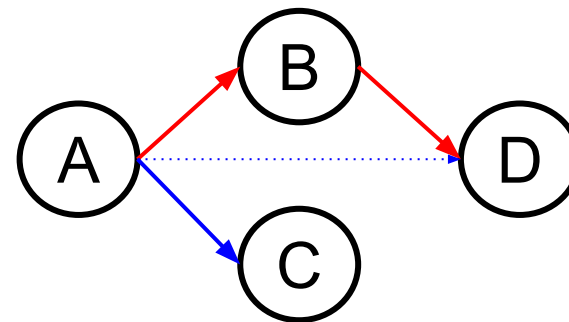
Read(d3)

```
d1' = Move d1
E' = compute new energy using d1', d2, d3
if metropolis_accept( $\Delta$ energy(E,E'), random_value())
  replace d1 with d1'
fi
```

## Background - Sequential task-flow (STF)

- The code is split into tasks
- Each task accesses the data in **READ** or **WRITE** (~**READ-WRITE**)
- **The tasks are created by a single thread** that informs the runtime system about the operations to do (which task to call with what data)
- The runtime system infers the dependencies between the tasks and ensures a correct execution (**sequential consistency**)
- DAG (node = task, edge = dependency)

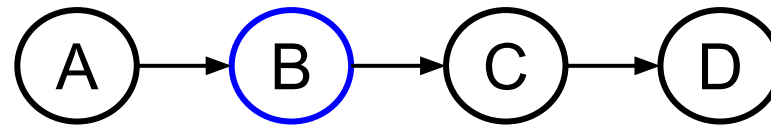
```
add task A: write(m), read(n) {  
    m += n;  
}  
add task B: write(n) {  
    n++;  
}  
add task C : write(l), read(m){  
    if(m > 0) l++;  
}  
add task D : read(m), read(n){  
    print(m, n);  
}
```



## TB Speculative execution example

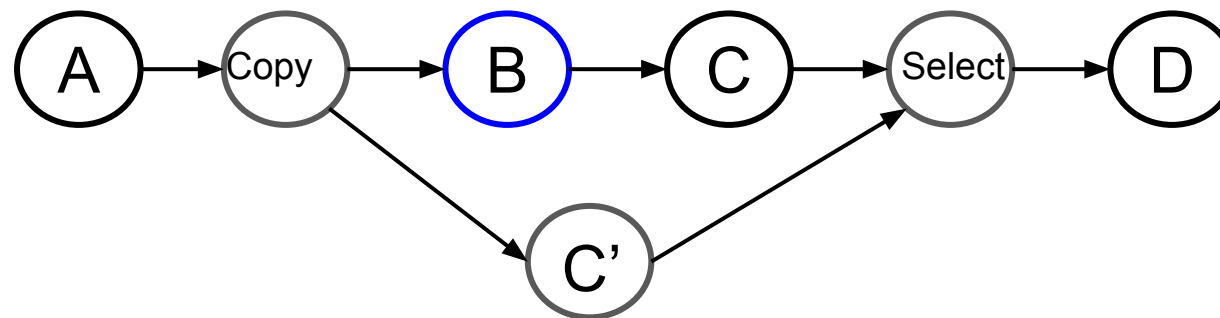
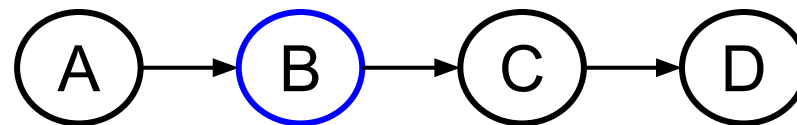
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- Consider here that a unique data is used by all tasks
  - Task B is an uncertain task (it is not sure if it will write on the data)



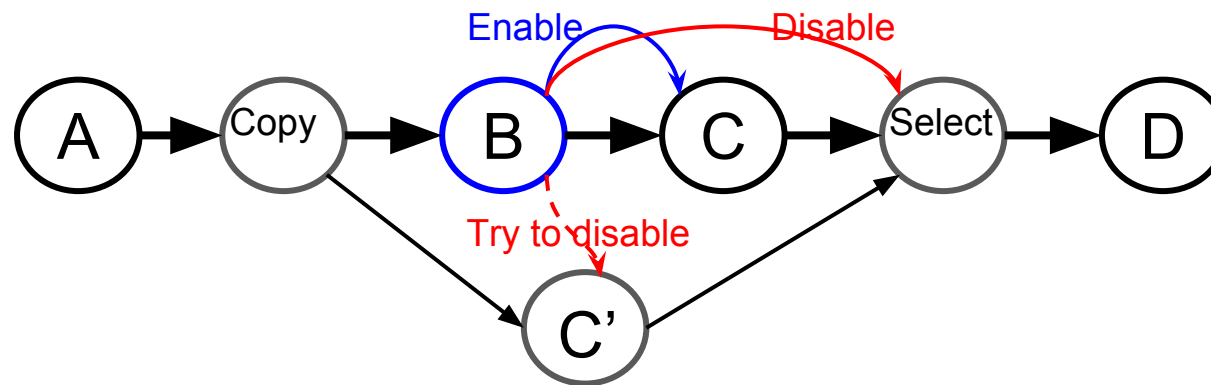
- Idea:
  - The programmer indicates that a task will potentially WRITE on a data
  - At the end of the task, the task informs the RS if the data was modified or not

## TB Speculative execution example

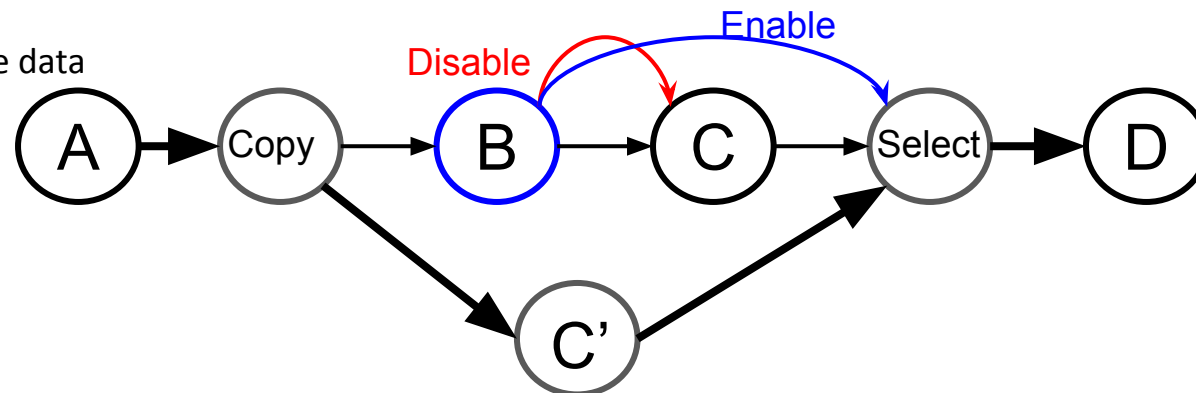


## TB Speculative execution example

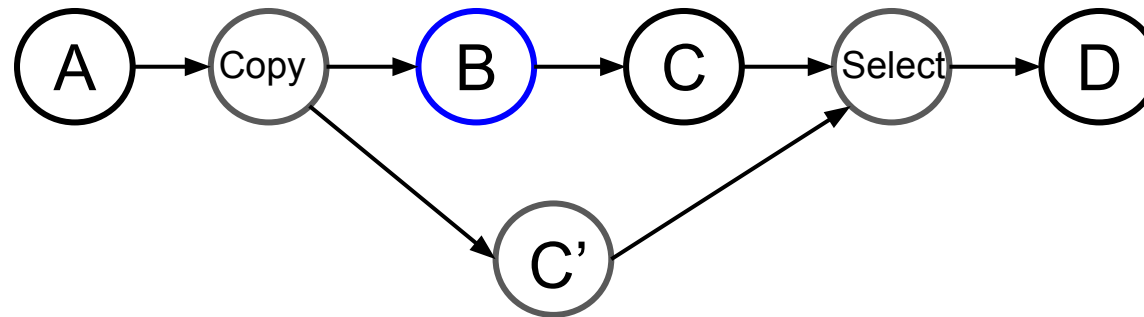
B wrote on the data



B did not write on the data



## TB Speculative execution example



- Execution time:
  - Without speculation :  $T(B) + T(C)$
  - With speculation :
    - If B writes on the data :  $T(\text{copy}) + \text{Max}(T(B) + T(C), T(C'))$  ( $T(C')$  is zero if canceled)
    - else :  $T(\text{copy}) + \text{Max}(T(B), T(C')) + T(\text{select})$

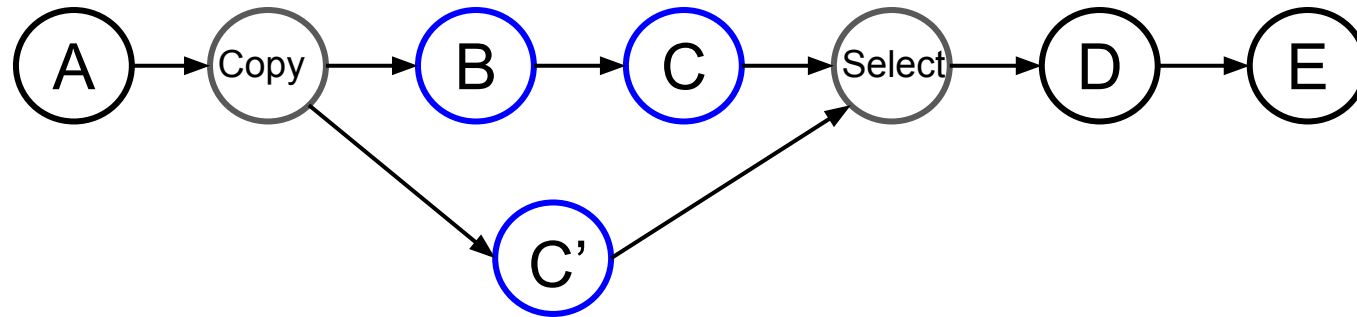
# Challenges

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- We do not have the full DAG from the beginning
  - Tasks are inserted dynamically (and executed)
  - We have to manage all the dependencies
- What to do when there are more than one uncertain task ? (partially solved)
- Unsolved/WIP
  - Schedule a DAG with speculation
  - Predict success/failure if speculation

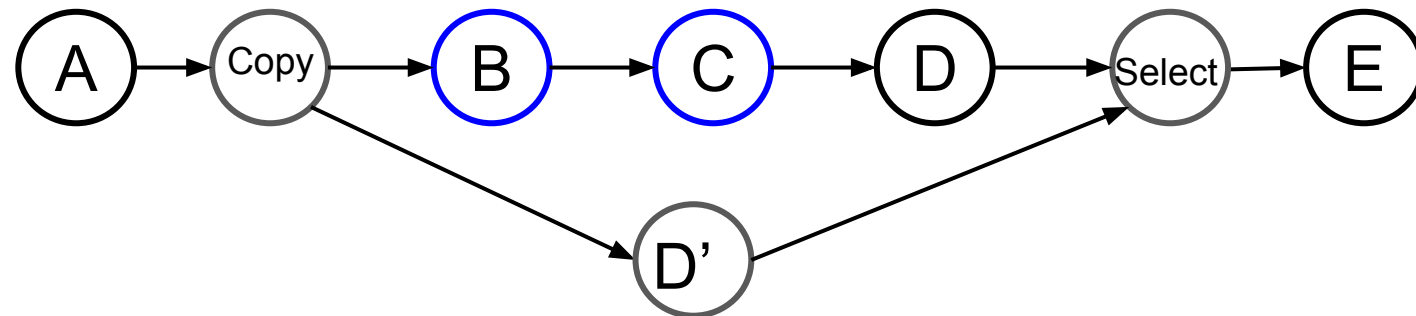
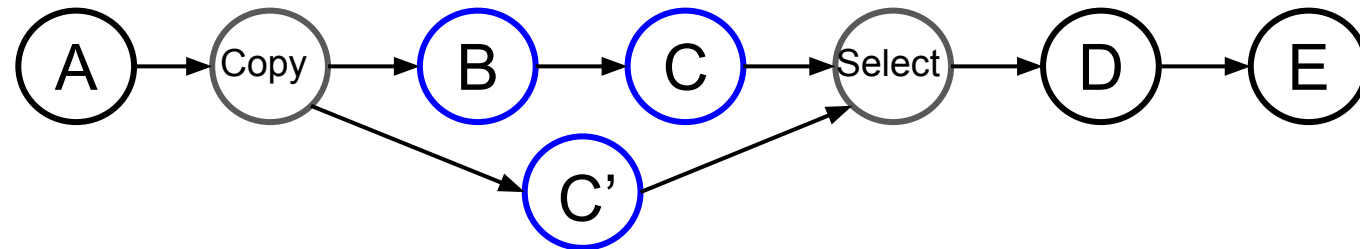
## TB Speculative execution example

- What if C is an uncertain task too?



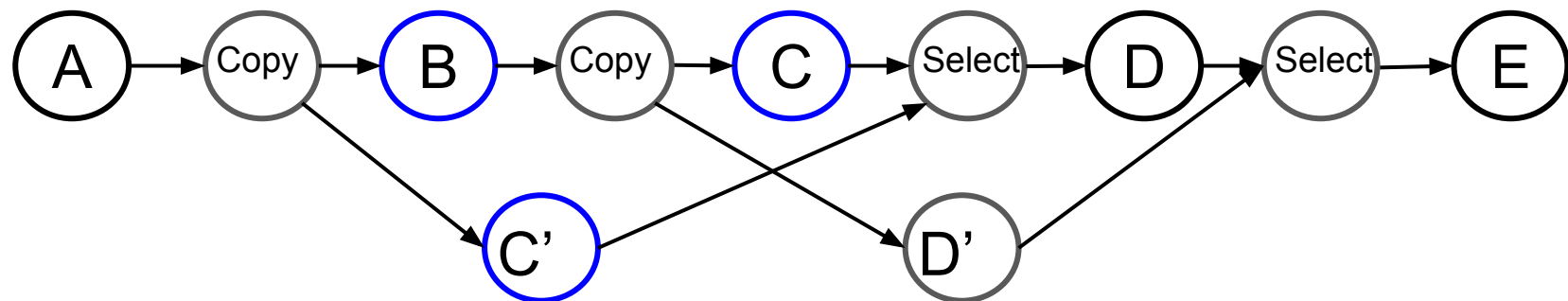
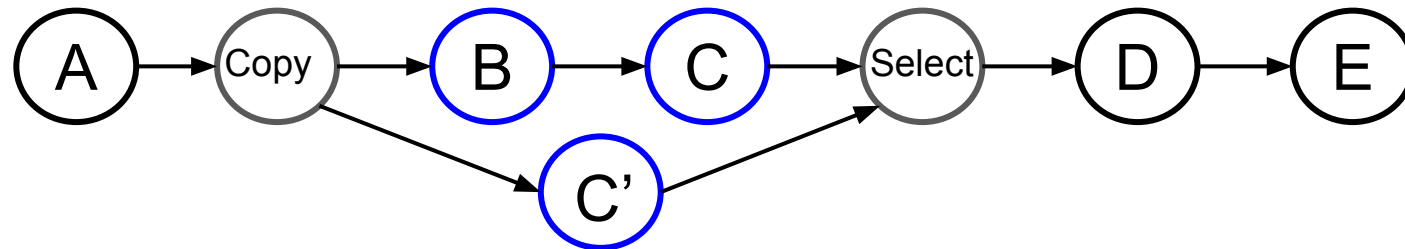


## TB Speculative execution example



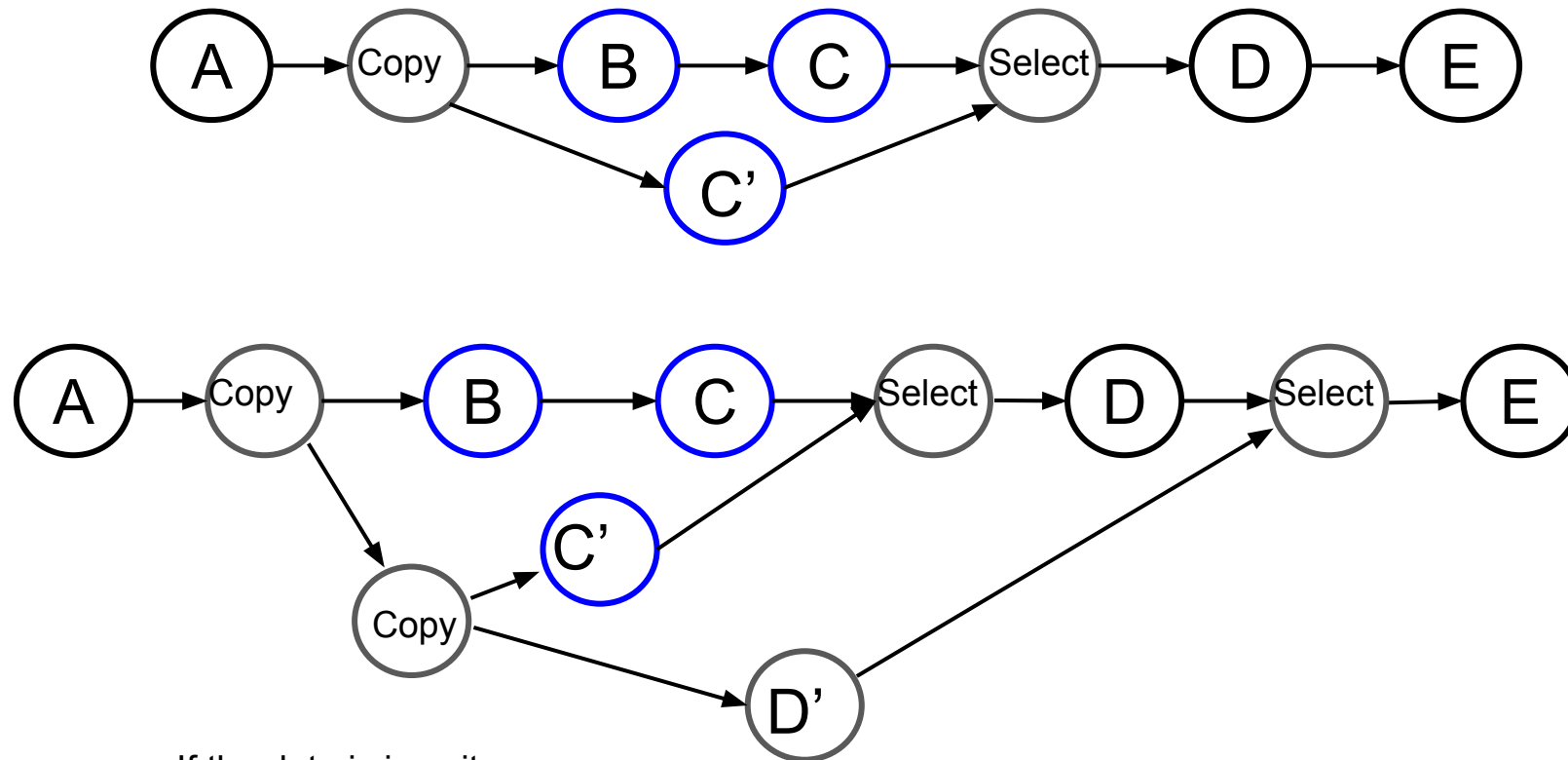
Poor parallelism, D is disabled only if both B and C did not modify the data, must compute  $B+C+(D \text{ or } D')$   
(This is not what we use)

## TB Speculative execution example



C', C and D' can be computed together  
(This is not what we use)

## TB Speculative execution example



If the data is in write  
 ((B,C) or C')D or D' can be computed together  
 (This is what we use)

# SPETABARU: SPeculative TAsk BAsed RUntime

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- A light C++ runtime (less than 3.000 lines, originally....)
  - Allow for compile time checking on the types
- Modern C++17
  - Advanced meta-programming
  - Lambda/anonymous functions
- Support data modes:
  - Read, Write, Atomic\_write, Commute
- Dynamic “array view”
  - To avoid ugly code when having dynamic dependencies

# SPETABARU: SPeculative TAsk BAsed RUnTime

- Example

```
// Create the runtime
const int NumThreads = SpUtils::DefaultNumThreads();
SpRuntime runtime(NumThreads);

const int initVal = 1;
int writeVal = 0;
// Create a task with lambda function
runtime.task(SpRead(initVal), SpWrite(writeVal),
    [](const int& initValParam, int& writeValParam){
        writeValParam += initValParam;
    });
// Create a task with lambda function (that returns a bool)
auto returnValue = runtime.task(SpRead(initVal),
    SpWrite(writeVal),
    [](const int& initValParam, int& writeValParam) -> bool {
        writeValParam += initValParam;
        return true;
    });
```

```
// Wait completion of a single task
returnValue.wait();
// Get the value of the task
const bool res = returnValue.getValue();
// Wait until two tasks (or less) remain
runtime.waitRemain(2);
// Wait for all tasks to be done
runtime.waitAllTasks();
// Save trace and .dot
runtime.generateTrace("/tmp/basis-trace.svg");
runtime.generateDot("/tmp/basis-dag.dot");
```

# SPETABARU: SPeculative TAsk BAsed RUnTime

---

```
SpRuntime runtime;  
  
runtime.setSpeculationTest([](const int /*inNbReadyTasks*/,  
                             const SpProbability& /*inProbability*/) -> bool {  
    // Always speculate  
    return true;  
});  
  
int val = 0;  
  
runtime.potentialTask(SpMaybeWrite(val), [](int& /*valParam*/) -> bool {  
    return false;  
}).setTaskName("C");
```

## SPETABARU: Example

Runtime.always\_speculate = true

Task A : read(val)

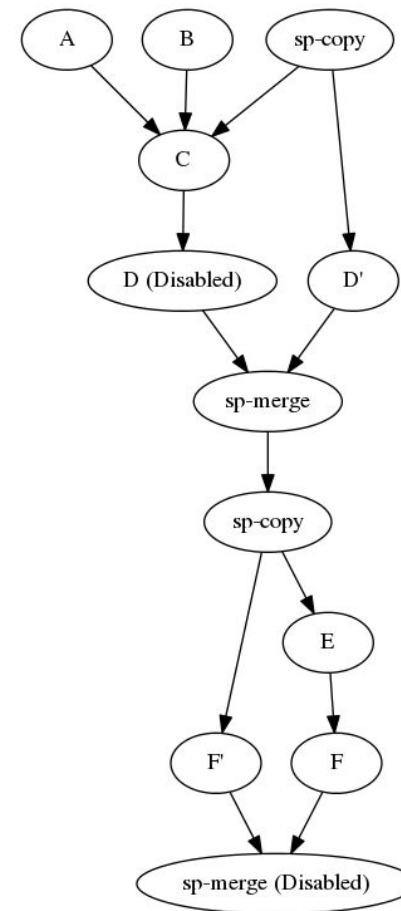
Task B : read(val)

Task C uncertain : maybe-write(val) -> false (no modifications)

Task D : write(val)

Task E uncertain : maybe-write(val) -> true (modifications)

Task F : write(val)



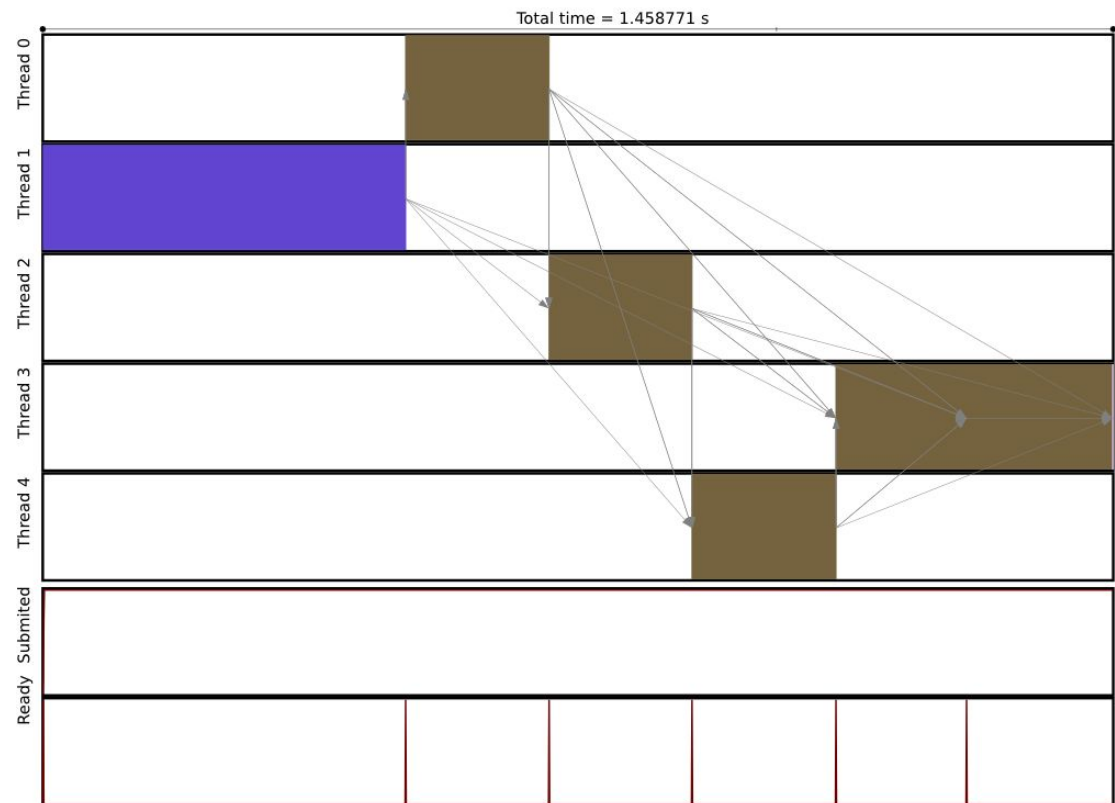
# Performance Study - MC

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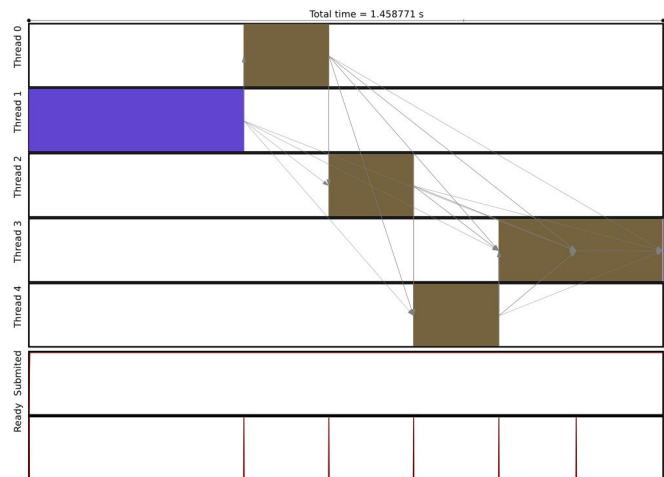
- Configuration
  - 5 domains
  - 2000 particles per domain
  - on a node with 32 cores but we use only 5
  - Reject/accept ratio  $\sim$  around 0.4 or 0.5
  - Number of iterations from 1 to 100
- 3 approaches:
  - Task-based (same as sequential)
  - Speculative (with speculation always enabled)
  - Speculative and force all move to be rejected (speculation performance limits, with speculation always enabled)



## Performance Study - MC (1 iteration - 5 threads)

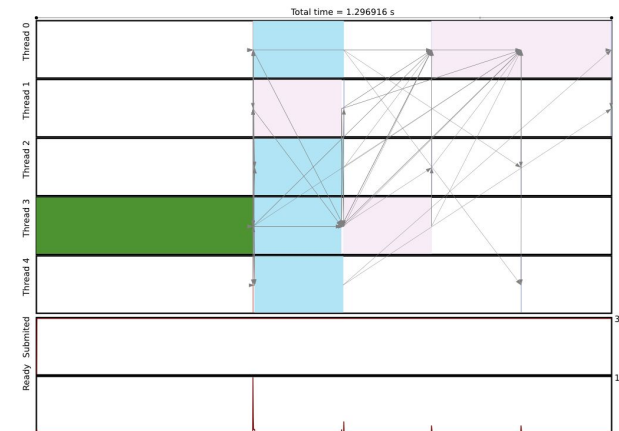


# Performance Study - MC (1 iteration - 5 threads)

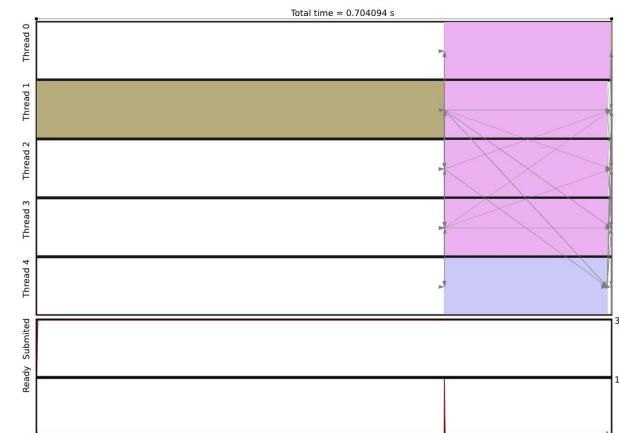


Without speculation - 1.45s

0 : reject  
**1 : accept**  
 2 : accepted  
 3 : reject  
 4 : reject

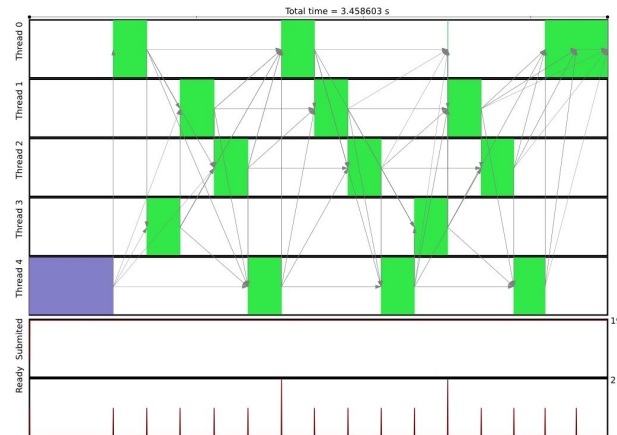


With speculation - 1.29s

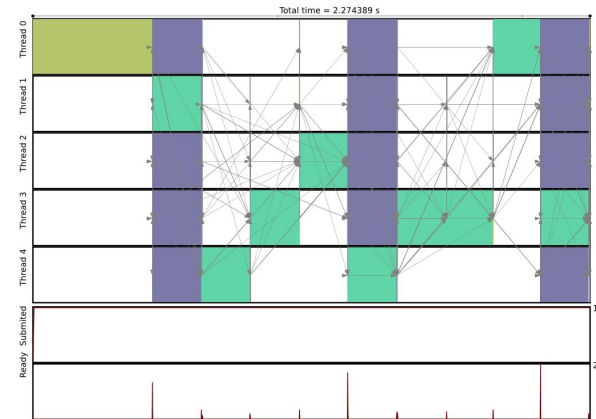


With speculation (all moves rejected) - 0.74s

# Performance Study - MC (3 iterations - 5 threads)



Without speculation - 3.45s

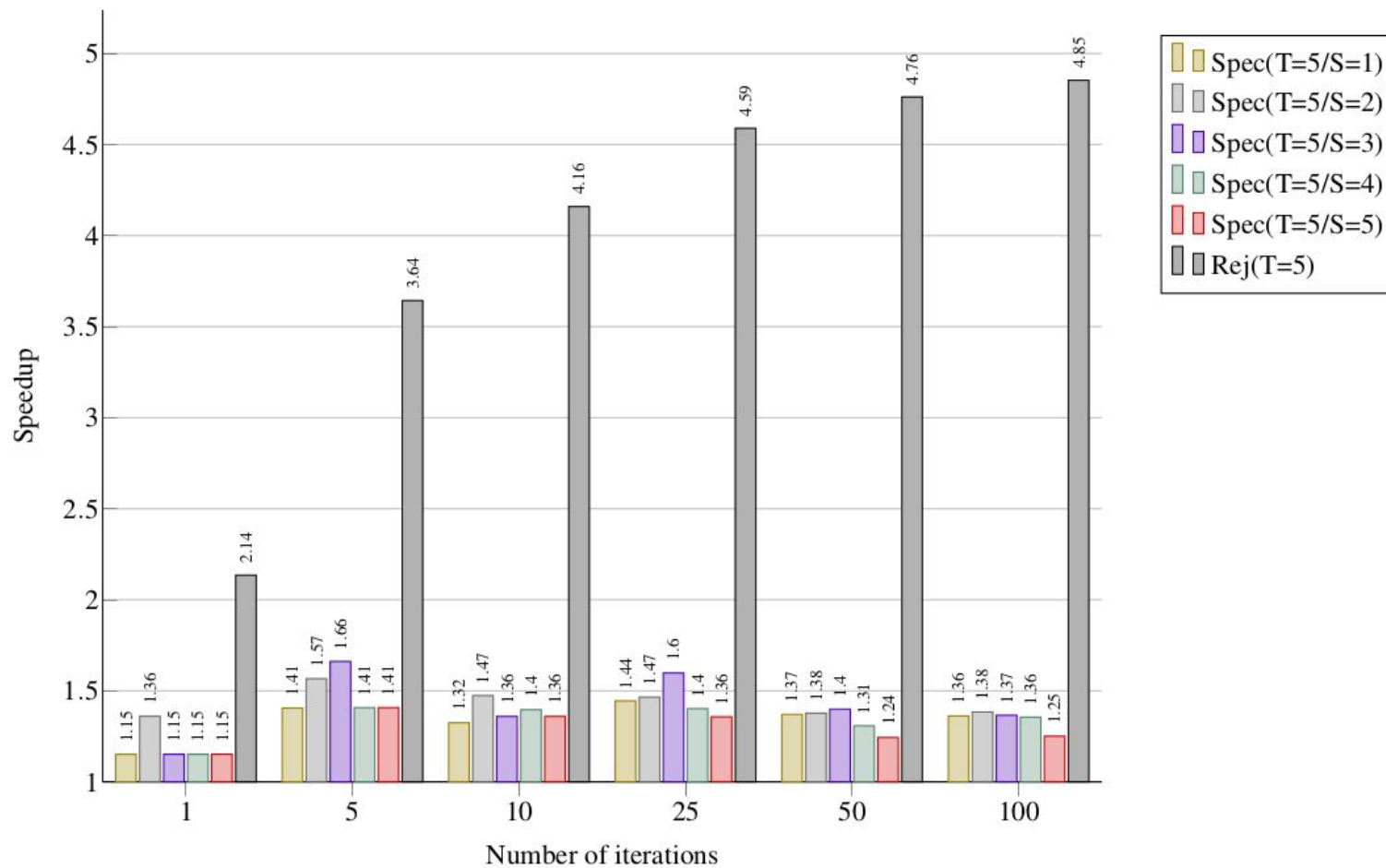


With speculation - 2.27s



With speculation (all moves rejected) - 1.09s

## Performance Study - MC - 5 threads

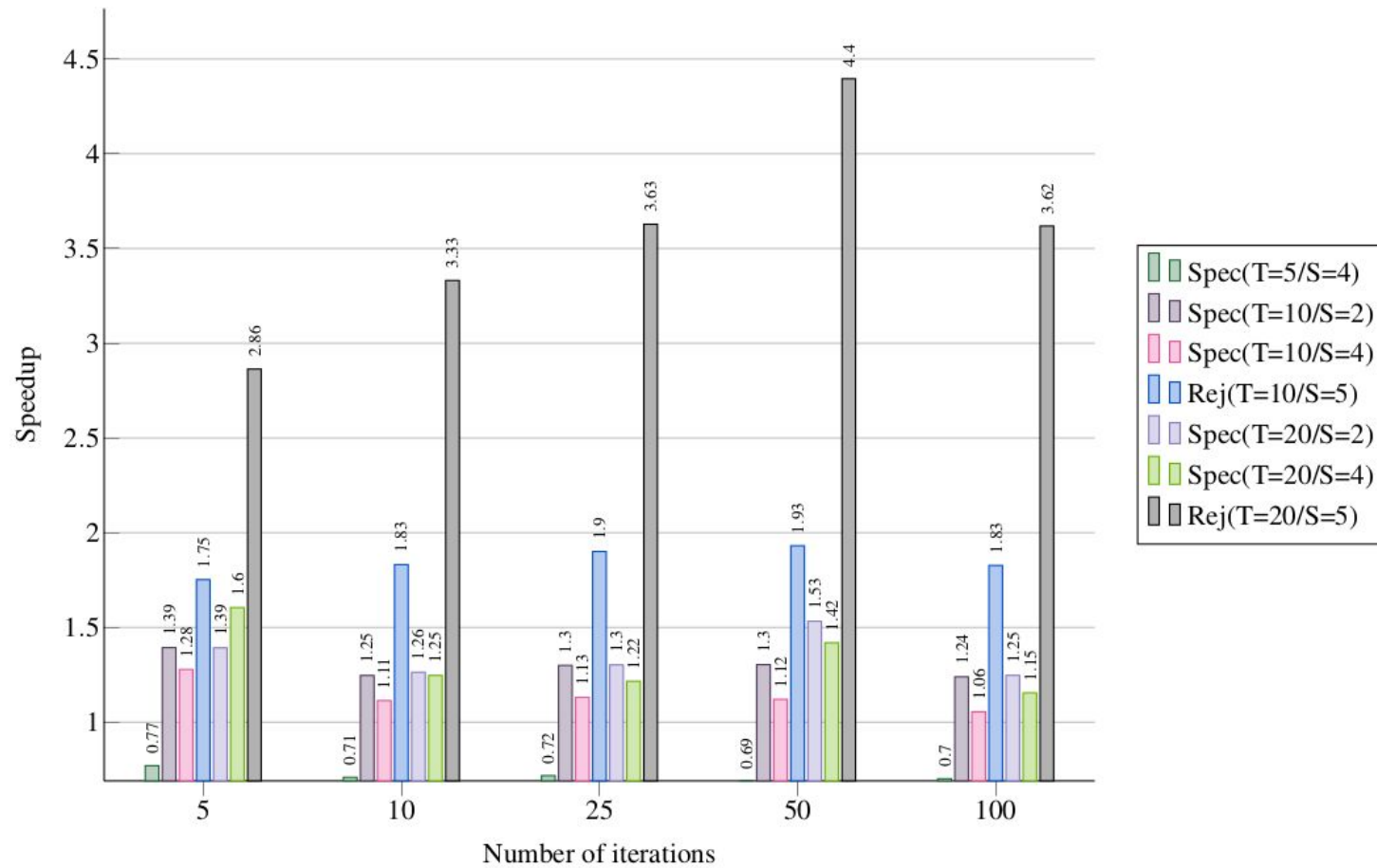


# Performance Study - Replica Exchange

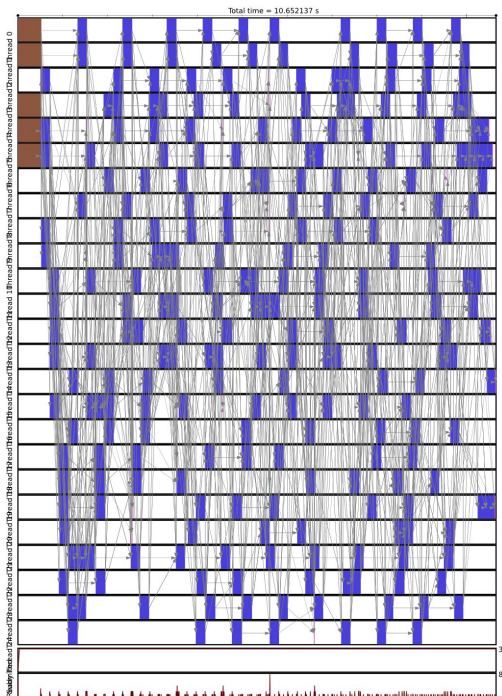
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- Configuration
  - 5 domains
  - 2000 particles per domain
  - on a node with 32 cores, we use 10 or 20 threads
  - Reject/accept rate  $\sim$  between 0.4 or 0.6
  - Number of iterations from 1 to 100
  - + 5 replicas
  - + Replica-exchange every 3 iterations
  - + Replica-exchange rate  $> 0.7$
- 3 approaches:
  - Task-based (same as one thread per replica)
  - Speculative (with speculation always enabled)
  - Speculative and force all move to be rejected (speculation performance limits, with speculation always enabled)

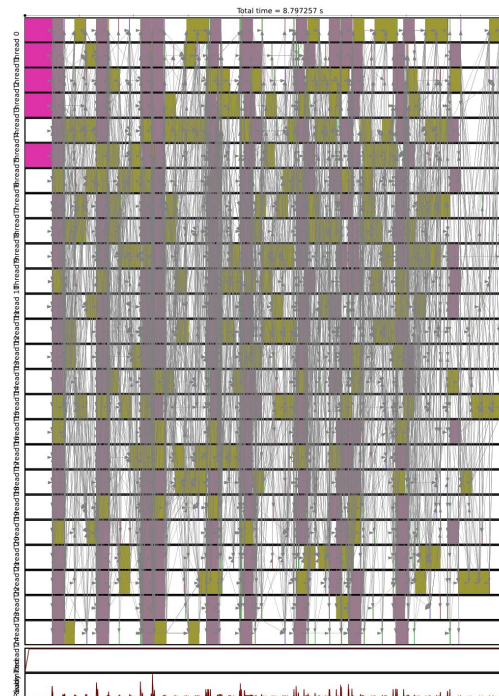
# Performance Study - Replica Exchange



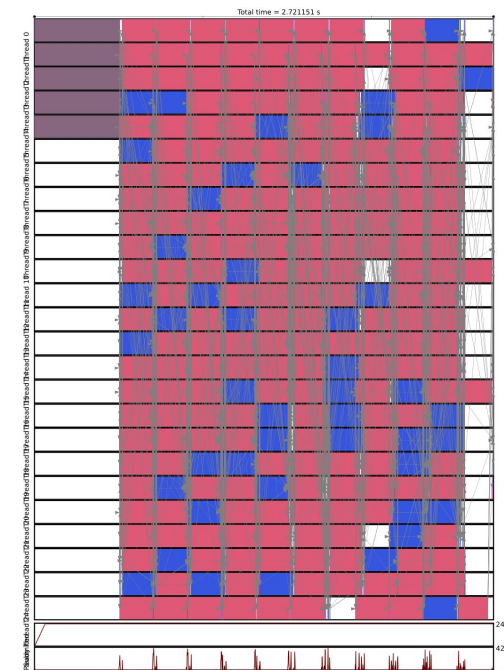
# Performance Study - Replica Exchange (10 iterations - 25 threads)



Without speculation - 10.6s



With speculation - 8.7s



With speculation (all moves rejected) - 2.7s

# Conclusions

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- First results to use speculation in TB method
- General pattern/algorithm
- SPETABARU is able to execute speculative task-flow
- Speedup for both the MC and REMC
  - obviously, the number of failures reduces the benefit
  - can slow down the executions if too much speculation (decision formula) and not enough threads



# Perspective

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- Limit the number of consecutive uncertain tasks automatically
- Inform RS if modifications have been done for each data individually (easily)
- Prediction of speculation success/failure with a decision formula
  - $\text{rand}() < \text{probability?}$
  - only if no tasks ready? seems not efficient
  - Use history with the information of the previous attempts
  - Use a perf. model
- Manage the scheduling of DAG with speculation
- Limit the number of copies (with a maximum memory occupancy value)
- Create other speculative execution models

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# Questions?

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[gitlab.inria.fr/bramas/spetabaru](https://gitlab.inria.fr/bramas/spetabaru)

